

Project leader: J.M. Le Goff
4 PhD students
<http://collspotting.web.cern.ch>



Leveraging insight into your data network by viewing co-occurrences while navigating across different perspectives.

Graphs

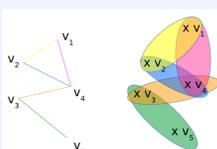
- Graph: set of vertices and set of edges.
- An edge links two vertices.
- Pairwise relationship.

Collaborations and (multi)sets

- Sets: group elements with no repetition and no order.
- Natural multisets: collection of objects with allowed repetitions.
- Co-occurrences are n -adic relationships.
- Co-occurrences are multisets, often reduced to sets.

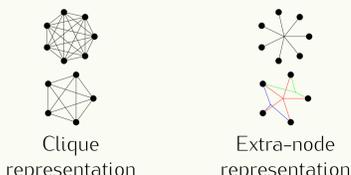
Hypergraphs

- Hypergraphs:
 - $\mathcal{H} = (V, E, i)$.
 - extend graphs.
 - allow relations between n vertices.
 - are a family of hyperedges E of nonempty subsets of vertices set V .
 - incident function is optional: $i: E \rightarrow \mathcal{P}(V)$.
- Hypergraphs are suited for co-occurrence modeling.

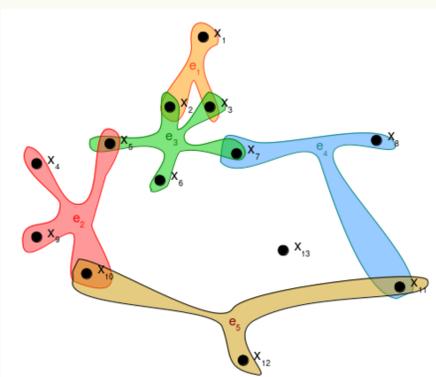


Hypergraph representations

Edge representation

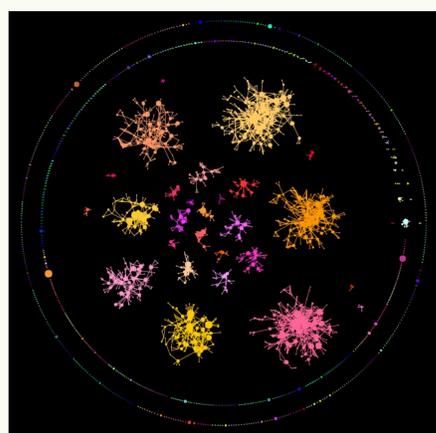


Set representation

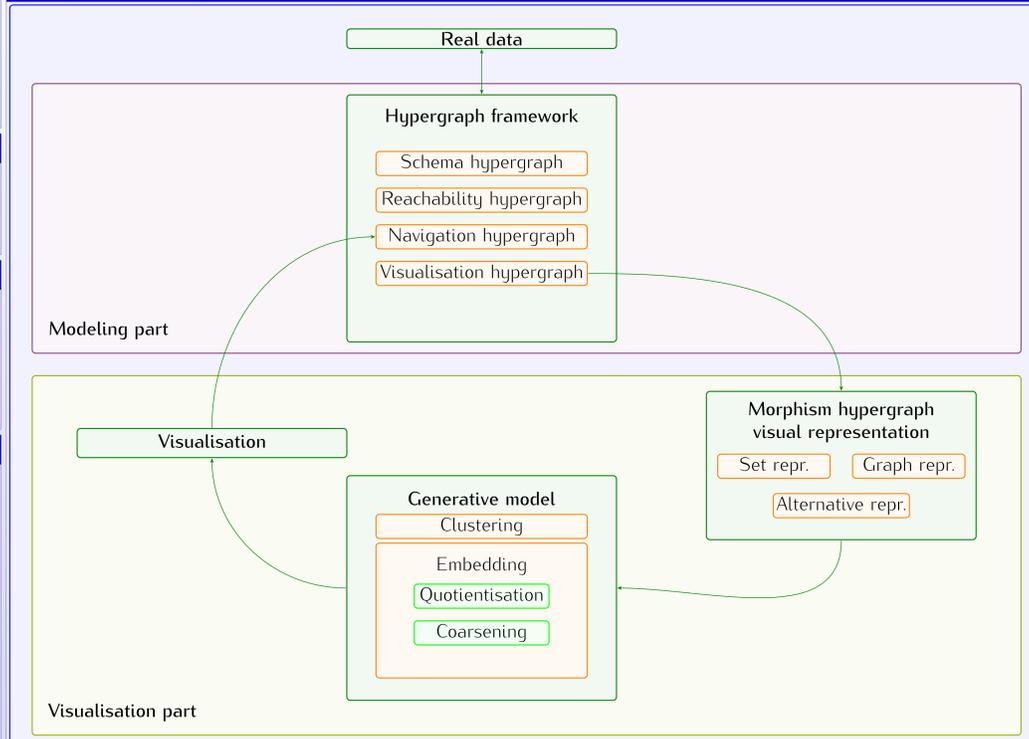


"PaintSplash" representation of hypergraph

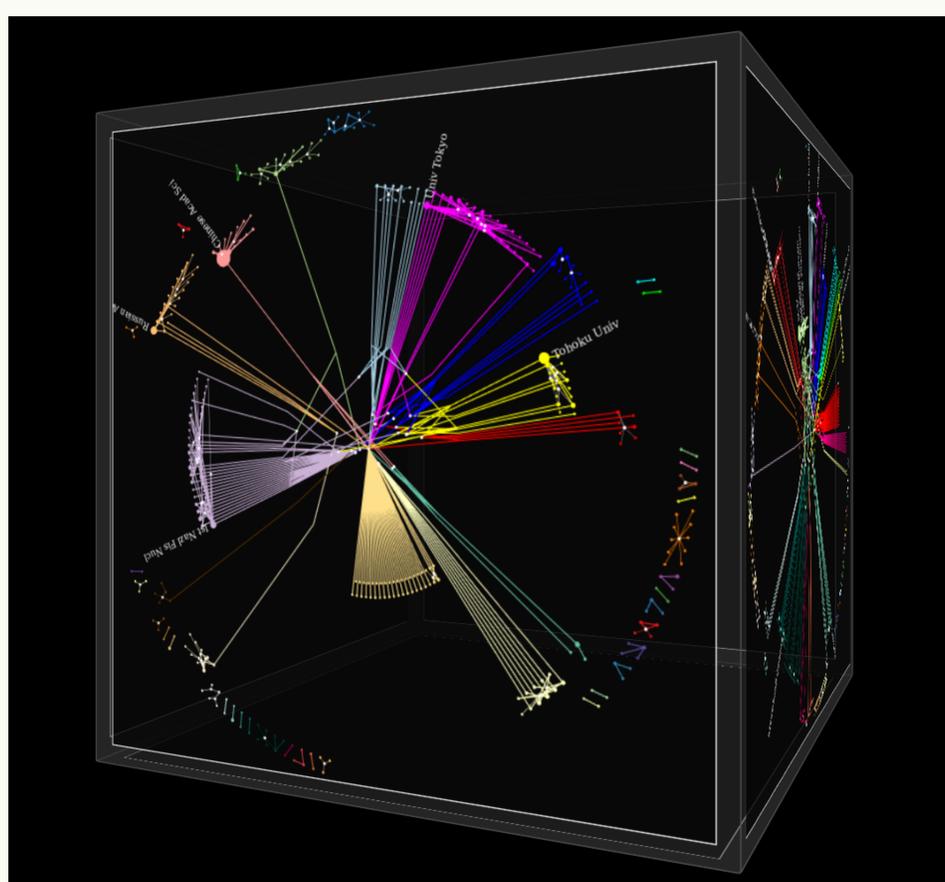
Large co-occurrence hypergraph



Hypergraph Modeling and Visualisation Framework



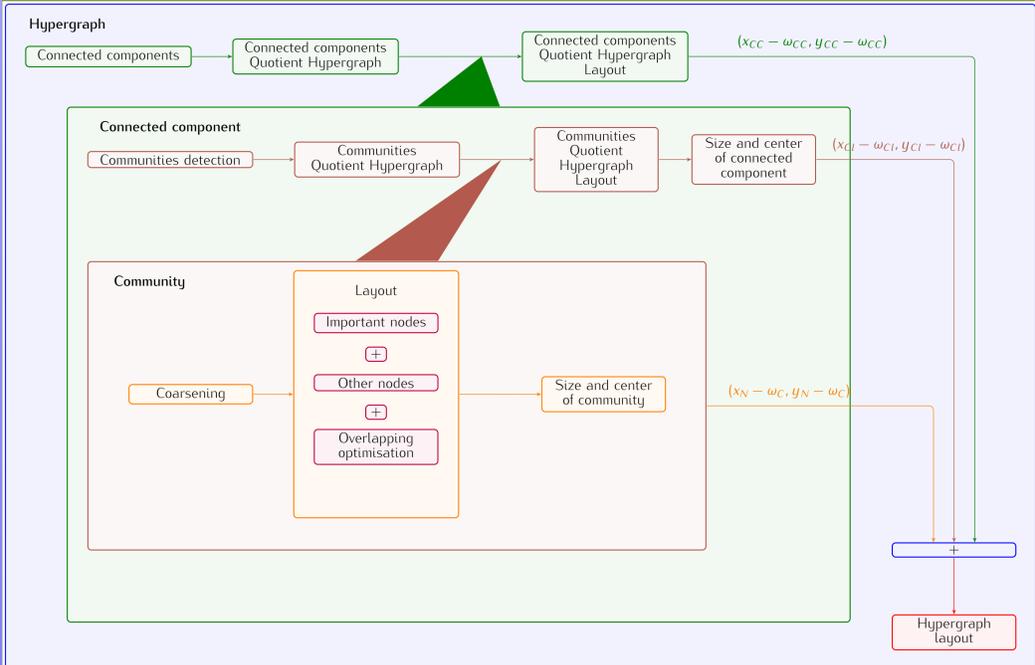
DataEdron



Improving visualisation in large hypergraphs

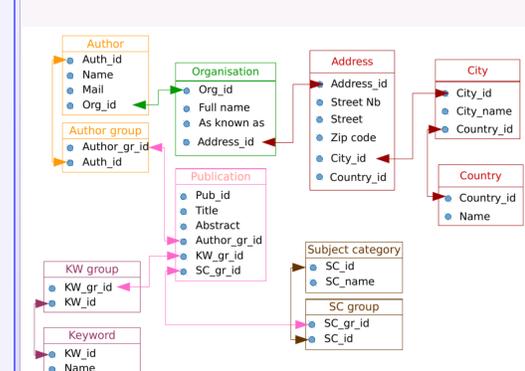
- Aim:**
 - keep meaningful information and structures;
 - focus on main information;
 - spreading the information uniformly on the layout.
- How can it be achieved?**
 - coarsening the visualisation hypergraph;
- What has been achieved?**
 - retrieval of important nodes using a diffusion process;
 - detection of important hyperedges of the network;
 - spectral comparison of the coarsened hypergraph with the original one.

Visualising large hypergraphs



Schema hypergraph

- In relational databases:
 - metadata instances => vertices.
 - tables => hyperedges.
 - foreign keys allow connection between hyperedges.
- In graph databases:
 - schema represents link between metadata.
 - schema not compulsory.
- => can be represented by a hypergraph.
- Schema hypergraph** $\mathcal{H}_{Sch} = (V_{Sch}, E_{Sch}, i_{Sch})$ represents the relations between metadata.



Schema hypergraph: exploded view. Shown on publication metadata example.

Extracted schema hypergraph \mathcal{H}_X

- only metadata instances of interest are kept in U , for:
 - visualisation;
 - reference;
 - search;
 - keeping connectivity in the hypergraph.
- $\mathcal{H}_X = (V_X, E_X, i_X)$, with $V_X = U$, $E_X = \{e \cap U : e \in E_{Sch}\}$ and $i_X = i_{Sch}|_{E_X}$.

Reachability hypergraph

- Reachability hypergraph** $\mathcal{H}_R = (V_R, E_R, i_R)$: obtained from \mathcal{H}_X by calculating its connected components.
- Vertices of \mathcal{H}_R : $V_R = V_X$.
- Hyperedges of \mathcal{H}_R : connected components of \mathcal{H}_X .

$$\forall e_R \in E_R : i_R(e_R) = \bigcup_{E_{CC} \subset e_R} \bigcup_{e \in E_{CC}} i_X(e)$$

Navigation hypergraph

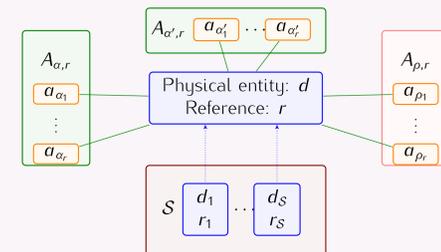
- Navigation hypergraph** $\mathcal{H}_N = (V_N, E_N)$: obtained from \mathcal{H}_R by choosing one hyperedge $e_R \in E_R$.
- Vertices of $\mathcal{H}_N = e_R$.
- Possible reference vertices in R_{ref} .
- Hyperedges of \mathcal{H}_N :

$$E_N = \{e_R \setminus R : R \subset R_{ref} \wedge R \neq \emptyset\}$$

- Navigation is possible without changing reference inside a hyperedge of \mathcal{H}_N .

Visualisation hypergraphs

- In a dataset \mathcal{D} , a physical entity d of reference r is fully described by: $(r, \{A_{\alpha,r} : \alpha \in V_{Sch}\})$.
- $A_{\alpha,r} = \{a_1, \dots, a_{\alpha}\}$: set of values of type α that are attached to d .



- For each $v \in \bigcup_{r \in S} A_{\rho,r} = \Sigma_{\rho}$, we build a set of physical references corresponding to data d that have v in attributes of type ρ : $R_v = \{r : v \in A_{\rho,r}\}$.
- Set of values of type α relatively to the reference v : $\bigcup_{r \in R_v} A_{\alpha,r} = e_{\alpha,v}$.
- Raw visualisation hypergraph** for the facet of type α/ρ attached to the search S is:

$$\mathcal{H}_{\alpha/\rho,S} = \left(\bigcup_{r \in S} A_{\alpha,r}, (e_{\alpha,v})_{v \in \Sigma_{\rho}} \right)$$

- By quotienting Σ_{ρ} and weighting => **reduced visualisation weighted hypergraph** for the search S :

$$\mathcal{H}_{\alpha/\rho,w,S} = \left(\bigcup_{r \in S} A_{\alpha,r}, \bar{E}_{\alpha}, W_{\alpha} \right)$$

Navigating through facets

- Reference type: ρ , current type α , target type: α' .
- Selecting vertices of type $A \subseteq A_{\alpha,S}$.
- Allows to:

- retrieve a subset of hyperedges of \bar{E}_{α} :

$$\bar{E}_{\alpha|_A} = \{e \in \bar{E}_{\alpha} \wedge (\exists x \in e : x \in A)\}$$

- retrieve the class \bar{v} attached to each $e \in \bar{E}_{\alpha|_A}$ => $\bar{V}|_A$ set of class \bar{v} .

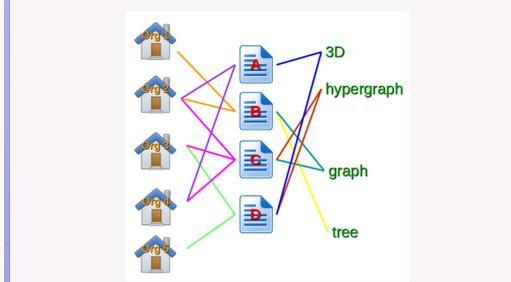
- retrieve the references of type ρ : $\mathcal{V}_{\rho,A} = \{v : \forall \bar{v} \in \bar{V}|_A : v \in \bar{v}\}$.

- R_v remains the same between facets => group of references: $S_A = \bigcup_{v \in \mathcal{V}_{\rho,A}} R_v$.

- switching to the facet of type α is then possible:

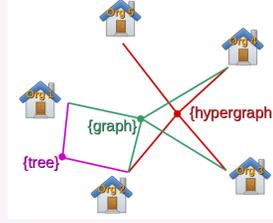
$$\mathcal{H}_{\alpha'/\rho|_A} = \left(\bigcup_{r \in S_A} A_{\alpha',r}, (e_{\alpha',v})_{v \in \mathcal{V}_{\rho,A}} \right)$$

Building co-occurrences on a publication dataset



Aim: Visualize co-occurrences of organisations in reference to keywords.

- Choose a type:**
 - α to visualize => organisations;
 - ρ to use as reference for co-occurrences => keywords.
- Result:** hypergraph with:
 - vertices: organisations;
 - hyperedges: organisation co-occurrences.



More info?

- Find my work on Arxiv: [1707.00115](https://arxiv.org/abs/1707.00115), [1712.08189](https://arxiv.org/abs/1712.08189), [1805.11952v2](https://arxiv.org/abs/1805.11952v2)
- Read more on diffusion by exchange: IEEE CBMI 2018 Proceedings.
- Read more on this: Proceedings of the 2nd IMA CTCDM.
- Contact: xavier.ouvrard@cern.ch